NEW TOAST IN TOWN

ONE TOAST FITS ALL

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Postgres breathed a second life into relational databases

• Postgres innovation - the first relational database with NoSQL support
• NoSQL Postgres attracts the NoSQL users
• JSON became a part of SQL Standard 2016
SQL/JSON in PostgreSQL

• The SQL/JSON construction functions: values of SQL types to JSON data

Mostly the same as json[b] construction functions

• JSON - generates JSON[b] from text data ( ::json[b] )
• JSON_SCALAR — generates a JSON[b] scalar value from SQL data (to_json[b])
• JSON_OBJECT - construct a JSON[b] object.
  • json[b]_build_object()
• JSON_ARRAY - construct a JSON[b] array.
  • json[b]_build_array()
• JSON_ARRAYAGG - aggregates values as JSON[b] array.
  • json[b]_agg()
• JSON_OBJECTAGG - aggregates name/value pairs as JSON[b] object.
  • json[b]_object_agg()
SQL/JSON: JSON

JSON - generates JSON[b] from text data (something like cast to json/jsonb types)

Syntax:

```
JSON (  
    expression [ FORMAT JSON [ ENCODING UTF8 ] ]  
    [ { WITH | WITHOUT } UNIQUE [ KEYS ] ]  
    [ RETURNING json_data_type ]
)
```
SQL/JSON: JSON

JSON - generates JSON[b] from text data (something like cast to json/jsonb types)

```sql
=# SELECT JSON('{{"a": 1, "a": 2} RETURNING JSON), JSON('{{"a": 1, "a": 2} RETURNING JSONB) as jsonb;

<table>
<thead>
<tr>
<th>json</th>
<th>jsonb</th>
</tr>
</thead>
<tbody>
<tr>
<td>{&quot;a&quot;: 1, &quot;a&quot;: 2}</td>
<td>{&quot;a&quot;: 2}</td>
</tr>
</tbody>
</table>

(1 row)

=# SELECT JSON('{{"a": 1, "a": 2} RETURNING JSON), JSON('{{"a": 1, "a": 2}'} WITH UNIQUE KEYS RETURNING jsonb;

ERROR: duplicate JSON object key value
```
SQL/JSON in PostgreSQL

• The SQL/JSON **retrieval** functions:

  • JSON_VALUE - Extract an SQL/JSON value from JSON data and return it as SQL scalar of specified type.
  • JSON_QUERY - Extract an SQL/JSON array or object from JSON data and returns JSON string
  • JSON_TABLE - Query a JSON text and present it as a relational table.
  • IS [NOT] JSON - test whether a string value is a valid JSON text.
  • JSON_EXISTS - test whether a JSON path expression returns any SQL/JSON items

• Supported only JSONB! Need GSON (generalized JSON API) to support JSON and JSONB without code complication.
SELECT x, JSON_VALUE(jsonb '{"a": 1, "b": 2}', '$.* ? (@ > $x)' PASSING x AS x
    RETURNING int
    DEFAULT -1 ON EMPTY
    DEFAULT -2 ON ERROR
) y
FROM generate_series(0, 2) x;

| x | y |
|---+---|
| 0 | -2 |
| 1 |  2 |
| 2 | -1 |

(3 rows)
### SQL/JSON examples: JSON_QUERY

```sql
SELECT
    JSON_QUERY(js::jsonb, '$'),
    JSON_QUERY(js::jsonb, '$' WITHOUT WRAPPER),
    JSON_QUERY(js::jsonb, '$' WITH CONDITIONAL WRAPPER),
    JSON_QUERY(js::jsonb, '$' WITH UNCONDITIONAL ARRAY WRAPPER),
    JSON_QUERY(js::jsonb, '$' WITH ARRAY WRAPPER)
FROM
    (VALUES
        ('null'),
        ('12.3'),
        ('true'),
        ('"aaa"'),
        (['1', null, "2" ]),
        ('{"a": 1, "b": [2]}')
    ) foo(js);
```

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<td>null</td>
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<td>&quot;aaa&quot;</td>
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<td>[1, null, &quot;2&quot;]</td>
</tr>
<tr>
<td>{&quot;a&quot;: 1, &quot;b&quot;: [2]}</td>
<td>{&quot;a&quot;: 1, &quot;b&quot;: [2]}</td>
<td>{&quot;a&quot;: 1, &quot;b&quot;: [2]}</td>
<td>{&quot;a&quot;: 1, &quot;b&quot;: [2]}</td>
<td>{&quot;a&quot;: 1, &quot;b&quot;: [2]}</td>
</tr>
</tbody>
</table>

(6 rows)
JSON_MODIFY
Missing SQL/JSON functionality
Example: add key "big" = true to all apartments having area greater than 70.

Simple desired jsonpath expression:

$.floor[*].apt[*]?(@.area > 70).big = true

Complex query with manual unnesting/aggregation:

SELECT jsonb_set(js, '{floor}', (SELECT jsonb_agg(jsonb_set(floor, '{apt}', (SELECT jsonb_agg(CASE WHEN jsonb_typeof(apt -> 'area') = 'number'
                                 AND (apt -> 'area')::int > 70
                                 THEN apt || '{"big": true}'
                                 ELSE apt END)
                          FROM jsonb_array_elements(floor->'apt') apt))
                   FROM jsonb_array_elements(js->'floor') floor))
FROM house;
Example: change contacts and street address.

Desired jsonpath expression:

\$\$.info.contacts = 'new contacts', \$.address.street = 'new address'

Ugly query with `jsonb_set()` chaining:

```
SELECT jsonb_set(jsonb_set(js, '{info,contacts}',
to_jsonb('new contacts'::text)), '{address,street}',
to_jsonb('new address'::text))
FROM house;
```

`UPDATE` can be done in more natural way with subscripting syntax:

```
UPDATE house SET js['address']['street'] = to_jsonb('new address'::text),
js['info']['contacts'] = to_jsonb('new contacts'::text);
```
JSON_MODIFY — syntax

```
JSON_MODIFY(jsonb_expr, operation, ...
    [RETURNING type]
    [PASSING expr AS name, ...])
```

Operations:

- SET jsonpath = expr [... ON EXISTING] [... ON MISSING] [... ON NULL]
- REPLACE jsonpath = expr [... ON MISSING] [... ON NULL]
- INSERT jsonpath = expr [... ON EXISTING] [... ON NULL]
- APPEND jsonpath = expr [... ON MISSING] [... ON NULL]
- REMOVE jsonpath [... ON MISSING]
- RENAME jsonpath WITH expr [... ON MISSING]
- KEEP jsonpath, ... [... ON MISSING] (not implemented yet !)
JSON_MODIFY – ON behaviors

ON NULL — executed when the new value is NULL
- NULL ON NULL — use JSON null as new value
- IGNORE ON NULL — do nothing
- ERROR ON NULL — raise an error
- REMOVE ON NULL — remove old value, if exists

ON EXISTING — executed when target JSON item exists
- IGNORE ON EXISTING — do nothing
- ERROR ON EXISTING — raise an error
- REPLACE ON EXISTING — replace old value with new value
- REMOVE ON EXISTING — remove old value

ON MISSING — executed when target JSON item does not exist
- IGNORE ON MISSING — do nothing
- ERROR ON MISSING — raise an error
- CREATE ON MISSING — insert new value
JSON_MODIFY – simplification example

Example: add key “big” = true to all apartments having area greater than 70.

Greatly simplified query using JSON_MODIFY:

```sql
SELECT JSON_MODIFY(js,
    SET '$.floor[*].apt[*] ? (@.area > $big_area).big' = true
    PASSING 70 AS big_area
) FROM house;
```

Complex query with manual unnesting/aggregation:

```sql
SELECT jsonb_set(js, '{floor}', (SELECT jsonb_agg(jsonb_set(floor, '{apt}', (SELECT jsonb_agg(CASE WHEN (apt -> 'area')::int > 70 THEN apt || '{"big": true}' ELSE apt END)
            FROM jsonb_array_elements(floor->'apt') apt))
    FROM jsonb_array_elements(js->'floor') floor)
FROM house;
```
Example: change contacts and street address.

Simplified query using JSON_MODIFY with multiple operations:

```sql
SELECT JSON_MODIFY(js, SET '$.info.contacts' = 'new contacts',
                   SET '$.address.street' = 'new address')
FROM house;
```

- There is no need to wrap text values into jsonb, SQL types are automatically mapped to corresponding SQL/JSON item types.
- Multiple operations can be executed in the single pass through jsonb, what can speed up execution by \(N_{\text{operation}}\) times (not yet implemented). This optimization is not possible in chained function calls.
CREATE TABLE test_object AS
SELECT i id, jsonb_build_object('x', i, 'z', (  
    SELECT jsonb_agg(x)
    FROM generate_series(1, (10.0 ^ (i / 10.0))::int) x)) jb
FROM generate_series(1,50) i;

SET $.x = 0
• JSON_MODIFY(jb, SET '$.x' = 0)
• jsonb_set(jb, '{x}', '0')
• jb || '{"x": 0}''
• (SELECT jsonb_object_agg(k, v)
    FROM (SELECT k, v FROM jsonb_each(jb) || kv(k,v) UNION SELECT 'x', '0') kv(k, v))

REMOVE $.x
• JSON_MODIFY(jb, REMOVE '$.x')
• jsonb_set_lax(jb, '{x}', NULL, false, 'delete_key')
• jb - 'x'
• (SELECT jsonb_object_agg(k, v)
    FROM jsonb_each(jb) kv(k,v) WHERE k <> 'x')

RENAME $.x
• JSON_MODIFY(jb, RENAME '$.x' WITH 'y')
• jsonb_set(jb - 'x', '{y}', jb -> 'x')
• (jb - 'x') || jsonb_build_object('y', jb -> 'x')
• (SELECT jsonb_object_agg(CASE k WHEN 'x' THEN 'y' ELSE k END, v)
    FROM jsonb_each(jb) kv(k, v))

PostgresPro
• Подробно о SQL/JSON смотри доклад на pgconf.ru-2022:

SQL/JSON Functions: Commit

- First discussed at developers meeting@Fosdem, Jan 28, 2017 in Brussels
- Post in -hackers, Feb 28, 2017, March CommitFest
- Committed March, 2022!

"...Nikita Glukhov (who probably deserves an award for perseverance)...

Hi there,

Attached patch is an implementation of SQL/JSON data model from SQL-2016 standard (ISO/IEC 9075-2:2016(E)), which was published 2016-12-15 and is available only for purchase from ISO web site (https://www.iso.org/standard/63556.html). Unfortunately I didn't find any public sources of the standard or any preview documents, but Oracle implementation of json support in 12c release 2 is very close (http://docs.oracle.com/database/122/ADJSN/json-in-oracle-database.htm), also we used https://livesql.oracle.com/ to understand some details.

Postgres has already two json data types – json and jsonb and implementing another json data type, which strictly conforms the standard, would be not a good idea. Moreover, SQL standard doesn't describe data type, but only data model, which "comprises SQL/JSON items and SQL/JSON sequences. The components of the SQL/JSON data model are:

1) An SQL/JSON item is defined recursively as any of the following:

a) An SQL/JSON scalar, defined as a non-null value of any of the following predefined (SQL) types:
Re: SQL/JSON features for v15

On 8/31/22 12:26 PM, Robert Haas wrote:
> On Wed, Aug 31, 2022 at 10:20 AM Jonathan S. Katz <jkatz(at)postgresql(dot)org> wrote:
>> Andres, Robert, Tom: With this recent work, have any of your opinions
>> changed on including SQL/JSON in v15?
>
> No. Nothing's been committed, and there's no time to review anything
> in detail, and there was never going to be.

OK. Based on this feedback, the RMT is going to request that this is reverted.

With RMT hat on -- Andrew can you please revert the patchset?
SQL/JSON Functions: Problems

SQL/JSON uses subtransactions for handling errors inside type casts and default expressions. This results to two problems:

1. ~2x slow down in the simplest cases due to subtransaction initialization overhead:

```
CREATE TABLE jsonb_nulls AS
SELECT 'null'::jsonb jb
FROM generate_series(1, 100000) i;

-- subtransaction is used for implicit NULL ON ERROR
SELECT JSON_VALUE(jb, '$') FROM jsonb_nulls;
Execution Time: 469.786 ms

-- no subtransaction, no error handling
SELECT JSON_VALUE(jb, '$') ERROR ON ERROR) FROM jsonb_nulls;
Execution Time: 225.519 ms
```
SQL/JSON Functions: Problems

2. Using functions that modify database in DEFAULT ON EMPTY leads to fast XID allocation per each processed row:

```sql
CREATE TABLE log (msg text);

CREATE FUNCTION log_msg (msg text) RETURNS text AS 'INSERT INTO log VALUES($1) RETURNING $1' LANGUAGE sql;

CREATE TABLE t AS
SELECT jsonb_build_object('a', i) jb FROM generate_series(1, 1000000) i; -- 1M of {"a": N}

SELECT txid_current(); => 745

-- ON EMPTY logs $.a and uses it as a result, it leads to error when $.a > 32767
SELECT JSON_VALUE(jb, '$.b' RETURNING int2 DEFAULT log_msg(jb->>'a') ON EMPTY) FROM t;

Execution Time: 16272.629 ms

-- 1M XIDs allocated (~50 000 XIDs per second)
SELECT txid_current(); => 1000747
SQL/JSON Functions: Problems

The same problem present in PL/pgSQL's BEGIN-EXCEPTION blocks which executed in subtransaction:

```sql
-- direct analogue of previous JSON_VALUE
CREATE OR REPLACE FUNCTION test(jb jsonb) RETURNS int2 AS $$
BEGIN
    IF NOT jb ? 'b' THEN
        RETURN log_msg(jb ->> 'a')::int2;
    END IF;
    RETURN (jb -> 'b')::int2;
EXCEPTION WHEN OTHERS THEN
    RETURN NULL;
END $$ LANGUAGE plpgsql;

SELECT txid_current(); => 1573

SELECT test(jb) FROM t;
Execution Time: 27584.998 ms

SELECT txid_current(); => 1001577
```
SQL/JSON Functions: Solutions

To remove subtransactions we can do the following:

- Completely disallow non-ERROR ON ERROR
- Limit the set of supported RETURNING types, support only types with known behavior of casts. Write special variants of cast functions with error passing instead of error throwing.

DEFAULT expressions:

- Disallow arbitrary DEFAULT ON EMPTY expressions, support only constant-like expressions that do not throw errors (constants, column references, parameters etc.)
- Consider function's volatility status (immutable, stable, volatile) for denying it in DEFAULT expressions and casts
- Disallow DEFAULT ON EMPTY
Why this talk?

- Startups want/need JSON[B]
- Blossom of Microservice architecture
- One-Type-Fits-All (JSON)
  - Client app — Frontend - Backend — Database
- JSONB is one of the main driver of Postgres popularity

Results of our experiments in 2021:

*Performance of JSONB (not only) can be improved by several orders of magnitude with proper modification of TOAST.*

- **Pluggable TOAST** - legal way to integrate our improvements into the Core
The Curse of TOAST: Unpredictable performance

CREATE TABLE test (jb jsonb);
ALTER TABLE test ALTER COLUMN jb SET STORAGE EXTERNAL;
INSERT INTO test
SELECT
    i id,
    jsonb_build_object(
        'id', i,
        'foo', (SELECT jsonb_agg(0)
            FROM generate_series(1, 1960/12))
    ) jb
FROM
    generate_series(1, 10000) i;

=# EXPLAIN(ANALYZE, BUFFERS) SELECT jb->'id' FROM test;

QUERY PLAN
-------------------------------------------------
Seq Scan on test   (actual rows=10000 loops=1)
    Buffers: shared hit=2500
    Planning Time: 0.050 ms
    Execution Time: 6.147 ms
(4 rows)

Small update cause significant slowdown!

=# UPDATE test SET jb = jb || '{"bar": "baz"}';
=# VACUUM FULL test; -- remove old versions

=# EXPLAIN (ANALYZE, BUFFERS) SELECT jb->'id' FROM test;

QUERY PLAN
-------------------------------------------------
Seq Scan on test (actual rows=10000 loops=1)
    Buffers: shared hit=30064
    Planning Time: 0.105 ms
    Execution Time: 38.719 ms
(4 rows)

Pageinspect: 64 pages with 157 tuples per page

WHY 30064 pages !!!!
TOAST Explained
The Oversized-Attribute Storage Technique

TOASTed (large field) values are compressed, then splitted into the fixed-size TOAST chunks (1996B for 8KB page)

• TOAST chunks (along with generated Oid chunk_id and sequence number chunk_seq) stored in special TOAST relation pg_toast.pg_toast_relid, created for each table with TOASTed attributes.

• TOASTed attribute in the original heap tuple is replaced with TOAST pointer (18 bytes) containing chunk_id, toast_relid, raw_size, compressed_size.

https://www.postgresql.org/docs/current/storage-toast.html
TOAST access

TOAST pointer refers (by Oid chunk_id) to heap tuples with chunks using B-tree index (chunk_id, chunk_seq). Overhead to read only a few bytes from the first chunk can be 3, 4 or even 5 index blocks.
The Curse of TOAST

Access to TOASTed JSONB requires reading at least 3 additional buffers:

• 2 TOAST index buffers (B-tree height is 2)
• 1 TOAST heap buffer
  • 2 chunks can be read from the same page, but if JSONB size > Page size (8Kb), then more TOAST heap buffers

```sql
=# EXPLAIN (ANALYZE, BUFFERS) SELECT jb->'id' FROM test;
QUERY PLAN
---------------------------------------------
Seq Scan on test (actual rows=10000 loops=1)
  Buffers: shared hit=30064
Planning Time: 0.105 ms
Execution Time: 38.719 ms
(4 rows)
```

<p>| | |</p>
<table>
<thead>
<tr>
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</thead>
<tbody>
<tr>
<td>Table</td>
<td>64</td>
</tr>
<tr>
<td>TOAST index</td>
<td>2 * 10000</td>
</tr>
<tr>
<td>TOAST table</td>
<td>1 * 10000</td>
</tr>
<tr>
<td><strong>Total</strong></td>
<td><strong>30064</strong></td>
</tr>
</tbody>
</table>
The Curse of TOAST

• TOAST is a very stable technology, which just works!
• TOAST has several hard-coded strategies to work with different data types, but it is "ancient" and knows nothing about jsonb, arrays, and other non-atomic data types.
• TOAST makes no attempt to take into account a workload, for example, append-only data.
• TOAST works only with binary BLOBs, when the TOASTed attribute is being updated, its chunks are simply fully copied. The consequences are:
  • TOAST storage is duplicated
  • WAL traffic is increased in comparison with updates of non-TOASTED attributes, since the whole TOASTed values is logged
  • As a result, performance is too low

• It’s time to improve it!
Motivational example (synthetic test)

A table with 100 jsonbs of different sizes (130B-13MB, compressed to 130B-247KB):

CREATE TABLE test_toast AS
SELECT
  i id,
  jsonb_build_object(
    'key1', i,
    'key2', (select jsonb_agg(0) from
generate_series(1, pow(10, 1 + 5.0 * i / 100.0)::int)),-- 10-100k elems
    'key3', i,
    'key4', (select jsonb_agg(0) from
generate_series(1, pow(10, 0 + 5.0 * i / 100.0)::int)) -- 1-10k elems
  ) jb
FROM generate_series(1, 100) i;

Each jsonb looks like: **key1, loooong key2[], key3, long key4[]**.

We measure execution time of operator ->(jsonb, text) for each row by repeating it 1000 times in the query:

```sql
SELECT jb -> 'keyN', jb -> 'keyN', ..., jb -> 'keyN' FROM test_toast WHERE id = ?;
```
Motivational example: SELECT

Key access time for TOASTed (raw size > 100 Kb) jsonbs linearly increase with jsonb size, regardless of key size and position.
Motivational example: UPDATE

Key update time for TOASTed (raw size > 100 Kb) jsonbs linearly increase with jsonb size, regardless of key size and position.

```
UPDATE test_toast SET jb = jsonb_set (jb, {keyN,...}, ?); COMMIT;
```
Jsonb deTOAST improvements goal

Ideal goal: no dependency on jsonb size and position

• Access/Update time ~ O(level) + O(key size)

• Original TOAST doesn’t use inline, only TOAST pointers are stored.
  Utilize inline (fast access) as much as possible:
  • Keep inline as much as possible uncompressed short fields and compressed medium-size fields
  • Keep compressed long fields in TOAST chunks separately for independent access and update.
Jsonb deTOAST improvements (root level)

• Partial (prefix) decompression - eliminates overhead of pglz decompression of the whole jsonb – FULL deTOAST and partial decompression:
  \[\text{Decompress}(\text{offset}) + \text{Detoast}(\text{jsonb compressed size}),\]
  offset depends on key position

• Sort jsonb object key by their length – good for short keys
  \[\text{Decompress}(\text{key_rank} \times \text{key size}) + \text{Detoast}(\text{jsonb compressed size}),\]
  offset depends on key size

• Partial deTOAST and partial decompression (deTOASTing iterator)
  \[\text{Decompress}(\text{key_rank} \times \text{key size}) + \text{Detoast}(\text{key_rank}_c \times \text{key size})\]
Jsonb deTOAST improvements

• Compress_fields – compress fields sorted by size until jsonb fits inline, fallback to Inline TOAST.
  0(1) – short keys
  Decompress(key size) – mid size keys

• Shared TOAST – compress fields sorted by size until jsonb fits inline, fallback to store compressed fields separately in chunks, fallback to Inline TOAST if inline overfilled by toast pointers (too many fields).
  • Access
    0(1) – short keys
    Decompress(key size) – mid size keys
    Decompress(key size) + Detoast(key size) – long keys
  • Update
    0(inline size) – short keys (inline size < 2KB)
    0(inline size) + 0(key size) – keys in chunks
    0(jsonb size) – inline TOAST
Jsonb deTOAST improvements

- Per-chunk compression – compress each chunk separately instead of slicing value compressed as a whole. This does not give any speedup and only enables random access needed for the next optimization.
- Sliced deTOAST – read and decompress only the necessary chunks to speedup access to elements of long TOASTed arrays or objects. Access time stops to depend from the element offset.
  - Without random access we need decompress all the data from the beginning of jsonb to the end of needed element, including the whole JEntry[] array:
    ```
    array: \text{Decompress}(N_{\text{elems}} * 4 + \text{element\_offset} + \text{element\_size}) + \\
    \text{Detoast} \ (N_{\text{elems}} * 4 + \text{element\_offset} + \text{element\_size}) \\
    \text{object: Decompress}(N_{\text{keys}} * 8 + \text{key\_value\_offset} + \text{key\_value\_size}) + \\
    \text{Detoast} \ (N_{\text{keys}} * 8 + \text{key\_value\_offset} + \text{key\_value\_size})
    ```
  - With random access we decompress only chunk containing JEntry and chunks containing element value:
    ```
    array: \text{Decompress}(\text{element\_index} * 4 \mod 2KB) + \text{Detoast}(1 \text{ chunk}) + \\
    \text{Decompress}(\text{element\_size} + \sim 2KB) + \text{Detoast}(\text{element\_size} + \sim 2KB) \\
    \text{object: Decompress}(N_{\text{keys}} * 8 + \text{key\_name\_offset} + \text{key\_name\_size}) + \text{Detoast}(1 \text{ chunk}) + \\
    \text{Decompress}(\text{key\_value\_size} + \sim 2KB) + \text{Detoast}(\text{key\_value\_size} + \sim 2KB)
    ```
Jsonb deTOAST improvements

• Direct TOAST – store inline chunk TID array together with TOAST pointer for elimination of index TOAST lookups.
  • Only ~330 (2KB/6B) chunks TIDs can fit into inline area.
  • A ~10x speedup for random access to individual chunks of values of <= 660 KB size.
  • But VACUUM FULL is broken, because we need to rewrite TIDs in TOAST pointers in the main table when chunks are moved inside the TOAST table (indirect access by index helped to avoid this).

• Compress TIDs – compress inline chunk TID array by the same technique that used in GIN indexes for compression of posting lists.
  • Increases number of inline TIDs up to 6x, so up to ~4MB of chunked data can be accessed directly from the TOAST pointer.
Jsonb deTOAST improvements

• In-place updates for TOASTed jsonb:
  • Store new element values, their offsets and lengths together with TOAST pointer (some kind of diff) instead of rewriting TOAST chunk chains, if element’s size and type not changed (in-place update) and new value fits into inline.
  • Old values replaced with new ones during deTOASTing.

• Update:
  • $O(\text{element size})$ – if in-place update and new value fits into inline
  • $O(\text{array size})$ – otherwise
Implementation details

• We implemented all optimizations inside TOAST (TOAST+) with the help of Generic JSON interface (GSON), which allows easy experimenting. Details are in "Understanding Jsonb Performance", "Scaling Jsonb" talks (links in References slide).

• It is hard to imagine that PG community will accept our optimizations, which touched a lot of PG Core.

• The only legal way is to go by "Postgres way" - make TOAST extensible (with backward compatibility).

• TOAST API would allow developers to develop new *toasters* without touching the Postgres core (more details in Part 2).

• We moved all optimizations from TOAST+ to *jsonb_toaster* (TOASTER).
• JSON, JSONB — two implementations of users functions, too many overlaps.

• Coming SQL standard (JSON 2.0) will specify JSON data type.
  • Oracle 20c already introduced JSON data type, which stored in binary format OSON
  • GSON provides compatibility to SQL Standard (JSON data type) and performance of JSONB - SET SQL_JSON = JSONB | JSON;

• More details in "JSON and JSONB Unification (GSON)" talk.
GSON: Unification of Jsonb and Json

- Generic JSON API allows reuse the code of jsonb user functions and operators for json data type.
- The old code of json functions and operators is almost completely removed.
- Json and jsonb user functions have only the separate entry in which input datums are wrapped into Json structures. But they share the same body by the calling generic subroutine.
- SQL/JSON uses these generic subroutines and generic jsonpath API.
- It is easy to add new features like partial decompression/detoasting or slicing, different storage formats (jsonbc, bson, oson,...).
- GSON allows to implement GIN opclass for indexing JSON.
GSON: Jsonx – optimized Jsonb for Toaster

- We removed JSON from GSON for simplicity, it may be added in any time.
- GSON allows us to create JSONX – an internal representation of JSONB in jsonb toaster, where we moved all our optimizations.
- JSONB remains untouched.
Pluggable TOAST introduces Custom TOAST Pointer with information about *toaster* and its data. In this example (B-tree index is omitted for simplicity) `jsonb_toaster` is used for toasting tuple with `id=4`, while other tuples are processed by *default toaster*, for example, tuple with `id=3`.
JSONB SELECT: TOASTER vs MASTER
JSONB UPDATE: TOASTER vs Master

UPDATE test_toast SET jb = jsonb_set (jb, {keyN,...}, ?); COMMIT;
• Access to short keys (key1, key3) improved by 1000X, no dependency on size and positions
• Access to array elements also improved, but require more optimizations.
SELECT JSONB: Aggregated statistics

TOASTER improves access to medium and long size (2 orders of magnitude) JSONB with minimal overhead for inline data.

```
SELECT jb -> 'keyN' FROM test_toast WHERE id = ?;
```
UPDATE Speedup: TOASTER vs Master

TOASTER greatly improves (up to 3 orders) the performance of update (inline, in-place) of JSONB!

UPDATE test_toast SET js = jsonb_test (jb, {keyN,...}, ?);
UPDATE JSONB: Aggregated statistics

TOASTER improves update of medium and long size (2 orders) JSONB with minimal overhead for inline data.

UPDATE test_toast SET js = jsonb_test (jb, {keyN,…}, ?);
Appendable Bytea Toaster

A table with 100 MB bytea (uncompressed):

CREATE TABLE test (data bytea);
ALTER TABLE test ALTER COLUMN data SET STORAGE EXTERNAL;
INSERT INTO test SELECT repeat('a', 100000000)::bytea data;

Append 1 byte to bytea (TOAST):  UPDATE test SET data = data || 'x'::bytea;

Update on test (actual time=1359.229..1359.232 rows=0 loops=1)
  Buffers: shared hit=238260 read=12663 dirtied=25189 written=33840
  -> Seq Scan on test (actual time=155.499..166.509 rows=1 loops=1)
    Buffers: shared hit=12665
Planning Time: 0.127 ms
Execution Time: 1382.959 ms  Table size doubled to 200 MB, 100 MB of WAL generated.

Append 1 byte to bytea (BYTEA TOASTER):

Update on test (actual time=0.060..0.061 rows=0 loops=1)
  Buffers: shared hit=2
  -> Seq Scan on test (actual time=0.017..0.020 rows=1 loops=1)
    Buffers: shared hit=1
Planning Time: 0.727 ms
Execution Time: 0.496 ms (2750x)  Table size remain 100 MB, 143 bytes of WAL generated.
Results – query execution time

UPDATE test SET data = data || repeat('a', append_size)::bytea;

- OLD + APPEND SIZE
- APPEND SIZE
Results – WAL traffic

UPDATE test SET data = data || repeat('a', append_size):bytea;

<table>
<thead>
<tr>
<th>append size, bytes</th>
</tr>
</thead>
<tbody>
<tr>
<td>10</td>
</tr>
<tr>
<td>100</td>
</tr>
<tr>
<td>1K</td>
</tr>
<tr>
<td>10K</td>
</tr>
<tr>
<td>100K</td>
</tr>
<tr>
<td>1M</td>
</tr>
</tbody>
</table>

\( \bigcirc \) OLD + APPEND SIZE

\( \bigcirc \) INLINE OLD + APPEND SIZE

PostgresPro
Conclusions

• TOAST optimizations based on knowledge of JSONB format resulted in a significant (several orders of magnitude) speedup of SELECT and UPDATE queries. Ideal goal: no dependency on jsonb size and position:
  + SELECT (need more research on nested data)
  + Inline (short keys), in-place (diff) UPDATE (need more work)

• TOAST can be optimized for specific workload, see for example, "TOAST for appendable bytea" talk (link in References slide), where we demonstrated streaming of binary data into Postgres, practically impossible for vanilla version.

• Postgres Extensibility allows us to Extend Postgres further:
  Data type aware TOAST — Pluggable TOAST!
PLUGGABLE TOAST

ONE TOAST FITS ALL
References


• Our experiments:
  • Understanding Jsonb performance
  • JSON and JSONB Unification (GSON)
  • Appendable Bytea TOAST
  • Pluggable TOAST at Commitfest https://commitfest.postgresql.org/38/3490/
  • Jsonb_toaster @Github (check License.txt):
    https://github.com/postgrespro/postgres/tree/jsonb_toaster
  • Meet us at GitHub https://github.com/postgrespro/postgres/tree/toasterapi_clean

• SQL/JSON in PostgreSQL
What is The Toaster?

- A set of methods implementing base operations for storage of oversized attributes using TOAST API.
- Aware of internal details of data it deals with
- Could be specific for column (different columns in table could use different Toasters)...
  - Or datatype (jsonb, bytea, etc.)
  - Or even workload (bytea specific append operation)
- Default toaster — compatible with previous versions
CREATE EXTENSION custom_toaster;

select * from pg_toaster;

<table>
<thead>
<tr>
<th>oid</th>
<th>tsrname</th>
<th>tsrhandler</th>
</tr>
</thead>
<tbody>
<tr>
<td>9864</td>
<td>deftoaster</td>
<td>default_toaster_handler</td>
</tr>
<tr>
<td>32772</td>
<td>custom_toaster</td>
<td>custom_toaster_handler</td>
</tr>
</tbody>
</table>

SQL – install custom Toaster:

CREATE FUNCTION custom_toaster_handler(internal)
RETURNS toaster_handler
AS 'MODULE_PATHNAME'
LANGUAGE C;

CREATE TOASTER custom_toaster HANDLER custom_toaster_handler;

Behind the scenes – create new Toaster:
## Plug In Custom Toaster...

SQL – attach created extension to table column:

```sql
CREATE TABLE tst1 (  
c1 text STORAGE plain,  
c2 text STORAGE external TOASTER custom_toaster,  
id int4
);

ALTER TABLE tst1 ALTER COLUMN c1 SET TOASTER custom_toaster;

=# \d+ tst1

| Column | Type   | Collation | Nullable | Default | Storage | Toaster       |... |
|--------+---------+-----------+----------+---------+----------+----------------+---|
| c1     | text    |           |          |         | plain    | deftoaster     |... |
| c2     | text    |           |          |         | external | custom_toaster |... |
| id     | integer |           |          |         | plain    |                |... |

Access method: heap
...as Column Attribute

New table column attribute – “atttoaster”. Used to insert data, but not suitable for update and fetch (as for Storage and Compression options):

```
=# \d pg_attribute
Table "pg_catalog.pg_attribute"
<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Collation</th>
<th>Nullable</th>
<th>Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>attrelid</td>
<td>oid</td>
<td></td>
<td>not null</td>
<td></td>
</tr>
<tr>
<td>attname</td>
<td>name</td>
<td></td>
<td>not null</td>
<td></td>
</tr>
<tr>
<td>attstorage</td>
<td>&quot;char&quot;</td>
<td></td>
<td>not null</td>
<td></td>
</tr>
<tr>
<td>atttoaster</td>
<td>oid</td>
<td></td>
<td>not null</td>
<td></td>
</tr>
<tr>
<td>attcompression</td>
<td>&quot;char&quot;</td>
<td></td>
<td>not null</td>
<td></td>
</tr>
</tbody>
</table>
```

- Data in table can be TOASTed with different Toasters
- Dropping Toaster is not allowed – all data must be re-toasted with the default
How TOAST Works

- Attribute stored in TOAST relation and replaced with Toast Pointer
- 4 strategies depending on storage type
- EXTENDED and MAIN are compressed before TOASTing — impossible to use optimizations
- EXTERNAL are TOASTed as-is — any knowledge of internal structure or workload could be used for optimizations, and could be replaced with custom structure keeping storage type and TOASTed data
How Pluggable TOAST Works

- Attribute is replaced with Custom Toast Pointer created by specific Toaster
- TOAST and storage details are hidden from AM behind the Toaster
- Custom Toast Pointer header contains only data necessary for Toaster lookup and required by the Executor
- Toaster can use any knowledge of data structure and workflow
- Same data could be TOASTed with different Toasters
Under The Hood Of Pluggable TOAST
Current TOAST

TOAST is slicing and storing an oversized value externally to Tuple
- Part of Heap AM (Core)
- Not extensible
- Same strategy for all datatypes
- Not effective for structured data (json) or data required special workflow (bytea)
Pluggable TOAST

TOAST is storing a value outside Tuple, with means and storage depending on Toaster

- Detached from Heap AM
- TOAST API
- Extensible
- Reference Toasters effective for heavy data types – json and bytea
Custom Toast Pointer

• Existing External (Default) Toast Pointer – varatt_external (=16 bytes header, 4 bytes alignment)

```c
typedef struct varatt_external {
    int32 va_rawsize;  /* Original data size (includes header) */
    uint32 va_extinfo; /* External saved size (without header) and */
                    /* compression method */
    Oid va_valueid;   /* Unique ID of value within TOAST table */
    Oid va_toastrelid; /* RelID of TOAST table containing it */
} varatt_external;
```

• Introduce new Custom Toast Pointer – varatt_custom (>=12 bytes header, 2 bytes alignment)

```c
typedef struct varatt_custom {
    uint32 va_toasterdatalen; /* total size of toast pointer, < BLCKSZ */
    uint32 va_rawsize;  /* Original data size (includes header) */
    uint32 va_toasterid; /* Toaster ID, actually Oid */
    char va_toasterdata[FLEXIBLE_ARRAY_MEMBER];  /* Custom toaster data */
} varatt_custom;
```
TOAST API Abstraction

Table Access Method

TOAST Helper – search Toaster by Toaster ID, call Toaster functions via API

TOAST API – unified extensible TOAST interface

Toaster with TOAST API implementation – Default toaster is built-in, Custom toasters are plugged in as extensions

Storage – Toasters could use any Storage and Access Methods to store data, independently of TAM
Toast Helper

- Toaster lookup functions
- input Toast Pointer type recognition
- column::atttoaster attribute check
- call Toaster

```c
extern TsrRoutine *GetTsrRoutine(Oid tsrhandler);
extern TsrRoutine *GetTsrRoutineByOid(Oid tsroid, bool noerror);
extern TsrRoutine *SearchTsrCache(Oid tsroid);
extern bool validateToaster(Oid toasteroid, Oid typeoid, char storage, char compression, Oid amoid, bool false_ok);
```
Backwards Compatibility

Default Toaster works with original Toast Pointers, and is enabled by default until changed by
ALTER TABLE table ALTER COLUMN column SET TOASTER toaster;

Keep EXTERNAL TOAST Pointer for compatibility:
if (VARATT_IS_EXTERNAL_ONDISK(attr))
{
    attr = PG_DETOAST_DATUM(PointerGetDatum(attr));
}
else if (VARATT_IS_CUSTOM(attr))
{
    Oid toasterid = VARATT_CUSTOM_GET_TOASTERID(attr);
    TsrRoutine *toaster = SearchTsrCache(toasterid);
    CustomToastRoutine *vtable = toaster->get_vtable(PointerGetDatum(attr));
    vtable-> ...
typedef struct TsrRoutine {
    NodeTag     type;
    /* interface functions */
    toast_init init;       // Initialization, storage creation, etc
    toast_function toast;  // Toast
    update_toast_function update_toast; // Update toasted value
    copy_toast_function copy_toast;       // Copy toasted value
    detoast_function detoast;  // Detoast
    del_toast_function deltoast;        // Delete toasted value
    get_vtable_function get_vtable;     // Return table of virtual functions
    toastervalidate_function toastervalidate; // Toaster validation
}
Toaster Validation

/* validate definition of a toaster Oid */
typedef bool (*toastervalidate_function)(Oid typeoid,
                           char storage, char compression,
                           Oid amoid, bool false_ok);

- Validate method is required for Toaster compatibility check depending on type, compression and storage options;
- Validation function must be aware of Storage and Compression because this is currently handled outside of Toaster;
- Validation could take in account that Toaster could be specific for datatype and workload
Default TOAST via TOAST API

pg_toaster table initial contents

{ oid => '9864', oid_symbol => 'DEFAULT_TOASTER_OID',
  descr => 'default toaster',
  tsrname => 'deftoaster', tsrhandler => 'default_toaster_handler' },

Default toaster handler

static bool
genericValidate (Oid typeoid, char storage, char compression,
  Oid amoid, bool false_ok)
{
  return true;
}

Datum
default_toaster_handler(PG_FUNCTION_ARGS)
{
  TsrRoutine *tsrroutine = makeNode(TsrRoutine);
  tsrroutine->init = genericToastInit;
  tsrroutine->toast = genericToast;
  tsrroutine->detoast = genericDetoast;
  tsrroutine->deltoast = genericDeleteToast;
  tsrroutine->get_vtable = NULL;
  tsrroutine->copy_toast = NULL;
  tsrroutine->update_toast = NULL;
  tsrroutine->toastervalidate = genericValidate;
  PG_RETURN_POINTER(tsrroutine);
}
Pimp My API – The Virtual Functions Table

- User-defined
- Provides extensibility

```c
void *
get_vtable()
```

A void pointer can be casted to any type – object, function or container

Yo dawg we heard you like APIs so we put an API in our API so you can extend Postgres while you extend Postgres
Toaster with User Functions – bytea Toaster

Code Sample: SQL

```sql
CREATE FUNCTION bytea_toaster_handler(internal)
RETURNS toaster_handler
AS 'MODULE_PATHNAME'
LANGUAGE C;
CREATE TOASTER bytea_toaster HANDLER bytea_toaster_handler;
```

Code Sample: C

```c
static void *
bytea_toaster_vtable(Datum toast_ptr)
{
    ByteaToastRoutine *routine = palloc0(sizeof(*routine));
    routine->magic = BYTEA_TOASTER_MAGIC;
    routine->append = bytea_toaster_append;
    return routine;
}

PG_FUNCTION_INFO_V1(bytea_toaster_handler);
Datum bytea_toaster_handler(PG_FUNCTION_ARGS)
{
    TsrRoutine *tsr = makeNode(TsrRoutine);
    tsr->init = bytea_toaster_init;
    tsr->toast = bytea_toaster_toast;
    tsr->deltoast = bytea_toaster_delete_toast;
    tsr->copy_toast = bytea_toaster_copy_toast;
    tsr->update_toast = bytea_toaster_update_toast;
    tsr->detoast = bytea_toaster_detoast;
    tsr->toastervalidate = bytea_toaster_validate;
    tsr->get_vtable = bytea_toaster_vtable;
    PG_RETURN_POINTER(tsr);
}
```
Overhead of TOAST API

Default Toaster

INSERT into test_toast_p VALUES (i, repeat('a', size));

UPDATE test_toast_p SET jb = repeat('b', pg_column_size(jb)+1) where id=i;

SELECT jb FROM test_toast_p WHERE id=i;
TOAST API Deliverables

- TOAST API (SQL syntax, core changes, patch set for PG15-16)
- Documentation on SQL Toaster Syntax (patch set for PG15-16)
- Default toaster – built-in (Core changes, included in TOAST API patch set);
- Dummy toaster as an example for developers – extension /contrib/dummy_toaster
- Bytea appendable toaster – extension /contrib/bytea_toaster
- Jsonb toaster – extension /contrib/jsonb_toaster
Conclusion

• The legal way to introduce powerful optimizations to internal data storage without changing or modifying Table AM;
• Extensible TOAST, pluggable Toasters;
• Toasters could use any Table Access Method and vice versa;
• Custom Toasters could be distributed under different licenses;
• Toasters could be extended not only to store, but to compress, encrypt and do other data transformations;
• Additional patch set – jsonb toaster as an extension, GSON API
Open Issues

• Drop Toaster – whole DB needs to be re-toasted if DROP TOASTER possible, otherwise TOASTed tuples won’t be restored;
• Move Storage and Compression options inside TOAST
• CREATE AS SELECT and INSERT AS SELECT operations re-toast data with Default Toaster no matter which Toaster is used in source table unless the target table is created and Toaster set before insertion
TODO

TOAST API
• Introduce iterators to Default Toaster
• Documentation for Toasters
• Benchmark TOAST API and Toasters
• Support Table AM and Compression in Toaster

JSONB
• Updates for long chunked arrays (rewrite only affected chunks)
• GSON update API (maybe updatable iterators)
• GSON optimization (to reduce overhead of wrapping)
• TOAST for long scalars (needs JsonbValue refactoring)
• Support for renames of TOASTed keys (need more complex old-vs-new value comparison in `toast_update()`)
• Goal – commit Pluggable Toast to PG16!
Addendum 1 – Code Samples
typedef void (*toast_init)(Relation rel, Datum reloptions, LOCKMODE lockmode, bool check, Oid OIDOldToast);
typedef struct varlena* (*toast_function) (Relation toast_rel, Oid toasterid, Datum value, Datum oldvalue, int max_inline_size, int options);
typedef struct varlena *((*update_toast_function) (Relation toast_rel, Oid toasterid, Datum newvalue, Datum oldvalue, int options));
typedef struct varlena *((*copy_toast_function) (Relation toast_rel, Oid toasterid, Datum newvalue, int options));
typedef struct varlena *((*detoast_function) (Datum toast_ptr, int offset, int length));
typedef void (*del_toast_function) (Datum value, bool is_speculative);
typedef void * (*get_vtable_function) (Datum toast_ptr);
typedef bool (*toastervalidate_function) (Oid typeid, char storage, char compression, Oid amoid, bool false_ok);
Default Toaster Implementation

Datum
default_toaster_handler(PG_FUNCTION_ARGS)
{
    TsrRoutine *tsrroutine = makeNode(TsrRoutine);

    tsrroutine->init = genericToastInit;
    tsrroutine->toast = genericToast;
    tsrroutine->detoast = genericDetoast;
    tsrroutine->deltoast = genericDeleteToast;
    tsrroutine->get_vtable = NULL;
    tsrroutine->copy_toast = NULL;
    tsrroutine->update_toast = NULL;
    tsrroutine->toastervalidate = genericValidate;

    PG_RETURN_POINTER(tsrroutine);
}
Toaster API with Virtual Table Example

```c
static Datum
bytea_toaster_append(Datum d1, Datum d2)
{
    ...
}
static void *
bytea_toaster_vtable(Datum toast_ptr)
{
    ByteaToastRoutine *routine = palloc0(sizeof(*routine));
    routine->magic = BYTEA_TOASTER_MAGIC;
    routine->append = bytea_toaster_append;
    return routine;
}
PG_FUNCTION_INFO_V1(bytea_toaster_handler);
Datum
bytea_toaster_handler(PG_FUNCTION_ARGS)
{
    TsrRoutine *tsr = makeNode(TsrRoutine);
    tsr->init = bytea_toaster_init;
    tsr->toast = bytea_toaster_toast;
    tsr->deltoast = bytea_toaster_delete_toast;
    tsr->copy_toast = bytea_toaster_copy_toast;
    tsr->update_toast = bytea_toaster_update_toast;
    tsr->detoast = bytea_toaster_detoast;
    tsr->toastervalidate = bytea_toaster_validate;
    tsr->get_vtable = bytea_toaster_vtable;
    PG_RETURN_POINTER(tsr);
}
```
Custom Toast Pointer Code

/* TOAST pointers are a subset of varattrib_1b with an identifying tag byte */
typedef struct
{
    uint8 va_header;  /* Always 0x80 or 0x01 */
    uint8 va_tag;    /* Type of datum */
    char va_data[FLEXIBLE_ARRAY_MEMBER];  /* Type-specific data */
} varattrib_1b_e;

/* varatt_custom is a subset of varattrib_1b_e */
typedef struct varatt_custom
{
    uint16 va_toasterdatalen;  /* total size of toast pointer, < BLCKSZ */
    uint32align16 va_rawsize;  /* Original data size (includes header) */
    uint32align16 va_toasterid; /* Toaster ID, actually Oid */
    char va_toasterdata[FLEXIBLE_ARRAY_MEMBER];  /* Custom toaster data */
} varatt_custom;

#define VARDATA_1B_E(PTR) (((varattrib_1b_e *) (PTR))->va_data)
#define VARATT_IS_EXTERNAL(PTR) VARATT_IS_1B_E(PTR)
#define VARATT_IS_CUSTOM(PTR) \
    (VARATT_IS_EXTERNAL(PTR) && VARTAG_EXTERNAL(PTR) == VARTAG_CUSTOM)